

Binary Search



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```
// Input: [0,0,....,0,1,....,1,1]
// Output: Index of last 0, or -1 if there is none
int search(vector<int> const& a) {
    // Invariant: l==-1 or a[l]=0
    int l=-1, r=a.size();
    while (r - l > 1) {
        int m = l + (r - l)/2;
        if (a[m] == 0)
            l = m;
        else
            r = m;
    }
    return l;
}
```



Implicit representation of array with predicates.

$$P(x) \wedge (x < y) \implies P(y)$$

Example: Find index of value x in array a_0, \dots, a_{n-1}



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$$P(i) = (a[i] > x)$$

$$P(-1) = 0$$



Make sure that the list starts with 0 and ends with 1.

Add $-\infty$ to the front and ∞ to the end to avoid any problems.

Example 1



N employees

K filing cabinets that contain a_0, \dots, a_{K-1} documents

Assign employee i to consecutive cabinets l_i, \dots, r_i

All cabinets need to be covered by exactly one employee

Minimize $\max_i \{a_{l_i} + a_{l_i+1} + \dots + a_{r_i-1}\}$



Reduction to a simpler problem using binary search:

What is the minimum cost?



Can it be done with cost x ?

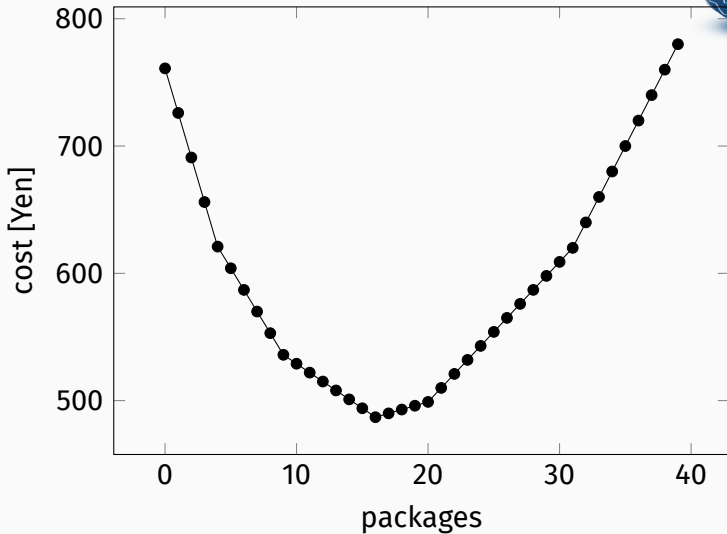
Example 2



Simple skewness of a collection of numbers: mean – median.

Given a list of n integers (not necessarily distinct), find the non-empty subsequence with the maximum simple skewness.

Search on Convex Function





- Binary search on derivative
- Ternary search (if you can't compute the derivative)

Example 3



Given a sequence of n integers a_1, a_2, \dots, a_n .

Determine a real number x such that the weakness of the sequence $a_1 - x, a_2 - x, \dots, a_n - x$ is as small as possible.

Weakness: maximum value of the poorness over all segments (contiguous subsequences) of a sequence.

Poorness of a segment: absolute value of the sum of the elements of segment.

Your answer will be considered correct if its relative or absolute error doesn't exceed 10^{-6}

Example 3: Take Away



```
// Input: [0,0,....,0,1,....,1,1]
// Output: Index of last 0
int search(double l, double r, double EPS) {
    // Invariant: a[l]=0
    while (r-l > EPS) {
        double m = (l + r)/2;
        if (pred(m))
            l = m;
        else
            r = m;
    }
    return l;
}
```



How many operations?



How many operations?

$$\mathcal{O}(\log((r - l)/e))$$



x : Solution

\tilde{x} : Approximation of x

Absolute error (difference to solution?):

$$|x - \tilde{x}|$$

Relative error (percentage of error?):

$$\frac{|x - \tilde{x}|}{|x|} = \left| \frac{\tilde{x}}{x} - 1 \right|$$



```
def IsApproximatelyEqual(x, y, epsilon=1e-6):  
    # Check absolute precision.  
    if -epsilon <= x - y <= epsilon:  
        return True  
  
    # Is x or y too close to zero?  
    if (-epsilon <= x <= epsilon or  
        -epsilon <= y <= epsilon):  
        return False  
  
    # Check relative precision.  
    return (-epsilon <= (x - y) / x <= epsilon or  
            -epsilon <= (x - y) / y <= epsilon)
```




Intuitive idea: Faster convergence for binary search on exponent

After some math:

$$m = \sqrt{lr} = e^{\frac{1}{2}(\log(l) + \log(r))}$$

Some pointer:

<http://codeforces.com/blog/entry/49189>



Doubles are discrete. They can be casted to 64 bit integers.

Find the best double to approximate a problem: Binary search!



```
int search(double l, double r) {
    for (;;) {
        uint64_t l_bits = *(uint64_t*)(&l);
        uint64_t r_bits = *(uint64_t*)(&r);
        if (r_bits - l_bits <= 1)
            return l;
        uint64_t m_bits = l_bits + ((r_bits - l_bits) >> 1);
        double m = *(double*)(&m_bits);
        if (pred(m))
            l = m;
        else
            r = m;
    }
}
```



Minimize absolute *and* relative error?

- Minimize absolute if $x < 1$.
- Minimize relative if $x > 1$.

Absolute error:

$$|x - \tilde{x}|$$

Relative error:

$$\frac{|x - \tilde{x}|}{|x|} = \left| \frac{\tilde{x}}{x} - 1 \right|$$



Needs to be larger than in task description because of rounding errors.

Can't be too large because the code will be too slow.

Hack: Fixed number of iterations!



Sometimes, there is no upper bound or you want your solution to be in $\mathcal{O}(\log(ans))$.

```
int r = 1;
while (!pred(r))
    r *= 2;
return search(l, r);
```



Can visualize binary search as a tree.

We look for the value that minimizes $(\text{Pred}(x), -x)$.

Branch and Bound: Define LB and UB => exactly the same as binary search.