

# Graph theory toolkit

SOI 2019, Sarnen, Gary Ye

February 11, 2019

# Introduction

- Previous lecture: Graph theory basics
- This lecture: Graph theory toolkit to solve many problems:
  - Implicit graphs
  - Reduction to graph problem
  - Modification of a graph problem
  - State space graph

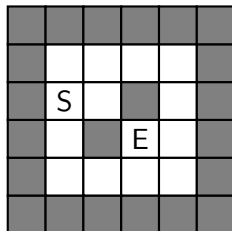


# Outline

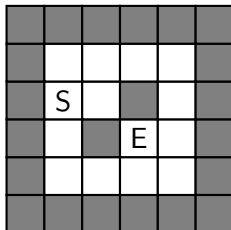
- 1 Implicit Graphs
- 2 Graph Duplication
- 3 State Space Graphs

# Problem: Maze

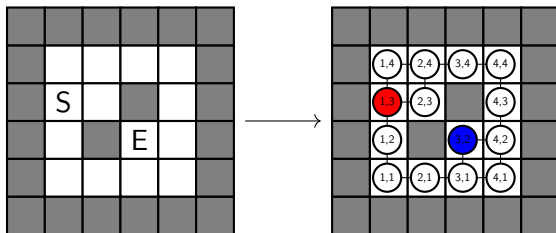
- Grey Cell ... Wall (unpassable)
- S ... Starting position
- E ... End position
- Each step: up, down, left or right.
- Goal: find path with minimum number of steps.



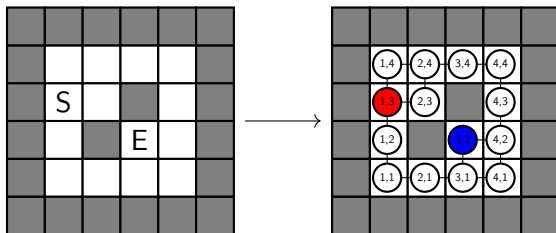
# Idea: Convert to Graph



# Idea: Convert to Graph

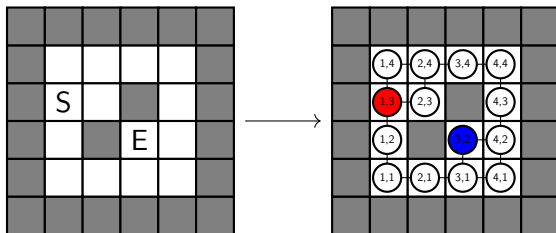


# Idea: Convert to Graph



→ Then apply shortest path algorithm on unweighted graph!

# Idea: Convert to Graph



- Then apply shortest path algorithm on unweighted graph!
- BFS!



# Graph representation: Explicit

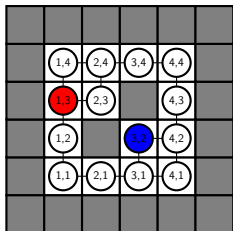
Store the graph explicitly in an adjacency list.

# Graph representation: Explicit

Store the graph explicitly in an adjacency list.

```
vector<vector<char>> maze = read(); // input
vector<vector<vector<pair<int,int>>> adj = init();
for(int i = 0; i < maze.size(); i++) {
    for(int j = 0; j < maze[i].size(); j++) {
        // code for adding edges to adj[i][j]
    }
}
```

## Graph representation: Explicit



Example:

$$\text{adj}[1][3] = \{(1, 4), \\ (2, 3), \\ (1, 2)\};$$

$$\text{adj}[3][2] = \{(4, 2), \\ (3, 1)\};$$

# Graph representation: Explicit

Downsides ☹:

- Additional memory required (up to  $\mathcal{O}(n^2)$  though)
- Additional implementation time required (only relevant for contests)

# Graph representation: Implicit

Instead of storing the edges explicitly, compute them using a simple rule:

## Graph representation: Implicit

Instead of storing the edges explicitly, compute them using a simple rule:

### Rule

There is an edge between cell  $c_1$  and cell  $c_2$  iff  $c_2$  is immediately adjacent to  $c_1$  and both cells do not represent a wall.

## Graph representation: Implicit

Instead of storing the edges explicitly, compute them using a simple rule:

### Rule

There is an edge between cell  $c_1$  and cell  $c_2$  iff  $c_2$  is immediately adjacent to  $c_1$  and both cells do not represent a wall.

### Explicit

```
// adj is precomputed
for(pair<int, int> q: adj[i][j]) {
    // do something with neighbor q
}
```

# Graph representation: Implicit

Instead of storing the edges explicitly, compute them using a simple rule:

## Rule

There is an edge between cell  $c_1$  and cell  $c_2$  iff  $c_2$  is immediately adjacent to  $c_1$  and both cells do not represent a wall.

## Explicit

```
// adj is precomputed
for(pair<int, int> q: adj[i][j]) {
    // do something with neighbor q
}
```

## Implicit

```
if(maze[i][j] != '#') {
    if(maze[i+1][j] != '#') {
        // do something with neighbor (i+1, j)
    }
    if(maze[i-1][j] != '#') {
        // do something with neighbor (i-1, j)
    }
    if(maze[i][j+1] != '#') {
        // do something with neighbor (i, j+1)
    }
    if(maze[i][j-1] != '#') {
        // do something with neighbor (i, j-1)
    }
}
```



# Maze: Lessons Learned

- We converted an ASCII grid problem into a graph problem
- We used an implicit graph representation to save memory and implementation time

# Outline

- 1 Implicit Graphs
- 2 Graph Duplication**
- 3 State Space Graphs

# Problem: Easteregg (SOI 2017)

## Formal statement

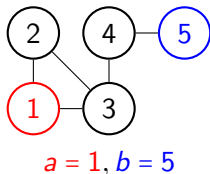
Given an undirected, unweighted graph with two starting vertices  $a$  and  $b$ . Choose a vertex  $c$  and two paths  $p = [a, \dots, c]$ ,  $q = [b, \dots, c]$  of equal length ( $|p| = |q|$ ) such that  $|p|$  (or  $|q|$ ) is minimized.

# Problem: Easteregg (SOI 2017)

## Formal statement

Given an undirected, unweighted graph with two starting vertices  $a$  and  $b$ . Choose a vertex  $c$  and two paths  $p = [a, \dots, c]$ ,  $q = [b, \dots, c]$  of equal length ( $|p| = |q|$ ) such that  $|p|$  (or  $|q|$ ) is minimized.

## Example

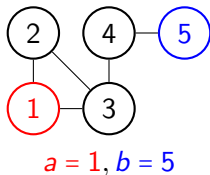


# Problem: Easteregg (SOI 2017)

## Formal statement

Given an undirected, unweighted graph with two starting vertices  $a$  and  $b$ . Choose a vertex  $c$  and two paths  $p = [a, \dots, c]$ ,  $q = [b, \dots, c]$  of equal length ( $|p| = |q|$ ) such that  $|p|$  (or  $|q|$ ) is minimized.

## Example



Solution: Choose  $c = 3$  and  $p = 1, 2, 3$  and  $q = 5, 4, 3$ .

# Problem: Easteregg (SOI 2017)

## Solution Idea

If  $P$  is the shortest path between  $a$  and  $b$ :

## Problem: Easteregg (SOI 2017)

### Solution Idea

If  $P$  is the shortest path between  $a$  and  $b$ :

→ choose the first half and the second half as the two paths and the vertex in the center vertex as the meeting point.

## Problem: Easteregg (SOI 2017)

### Solution Idea

If  $P$  is the shortest path between  $a$  and  $b$ :

→ choose the first half and the second half as the two paths and the vertex in the center vertex as the meeting point.

→ but only works if the length of the path  $P$  is even!



# Problem: Easteregg (SOI 2017)

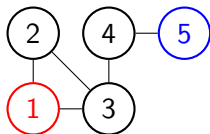
## Solution Idea

If  $P$  is the shortest path between  $a$  and  $b$ :

→ choose the first half and the second half as the two paths and the vertex in the center vertex as the meeting point.

→ but only works if the length of the path  $P$  is even!

## Illustration



# Problem: Easteregg (SOI 2017)

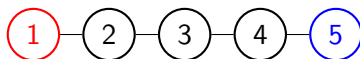
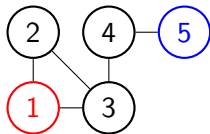
## Solution Idea

If  $P$  is the shortest path between  $a$  and  $b$ :

→ choose the first half and the second half as the two paths and the vertex in the center vertex as the meeting point.

→ but only works if the length of the path  $P$  is even!

## Illustration



The path above represents the shortest path with an even length, therefore we can choose  $c = 3$ .

## Problem: Easteregg (SOI 2017)

We have reduced our problem to an easier problem 😊:

Easier problem

Find the shortest path  $P$  between  $a$  and  $b$  of *even* length.

# Problem: Easteregg (SOI 2017)

We have reduced our problem to an easier problem 😊:

## Easier problem

Find the shortest path  $P$  between  $a$  and  $b$  of *even* length.

However ...

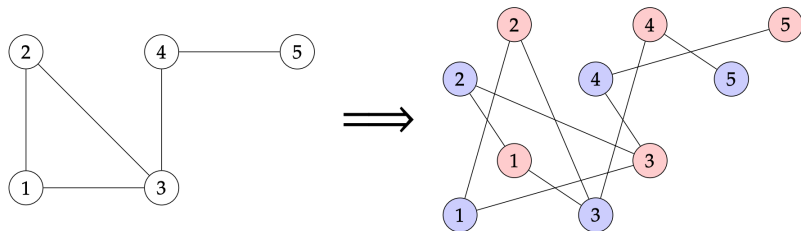
- Using BFS does not guarantee us that the path is even!
- If the shortest path  $P$  between  $a$  and  $b$  has an odd length, it does not mean that there exists no answer!

# Trick: Graph Duplication

We modify our graph by duplicating every node and every edge:

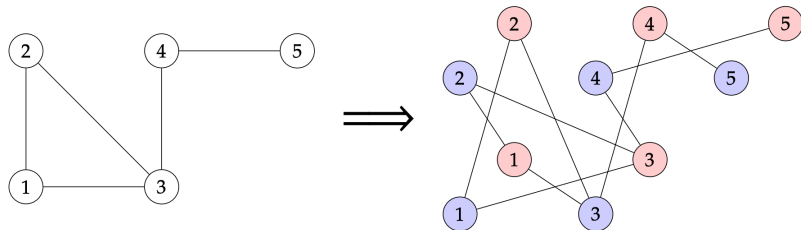
# Trick: Graph Duplication

We modify our graph by duplicating every node and every edge:



# Trick: Graph Duplication

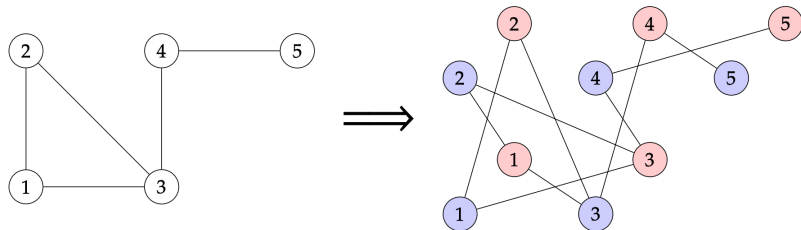
We modify our graph by duplicating every node and every edge:



→ whether the length of a path  $P$  is even can be checked by comparing the colors of the start and end point

# Trick: Graph Duplication

We modify our graph by duplicating every node and every edge:



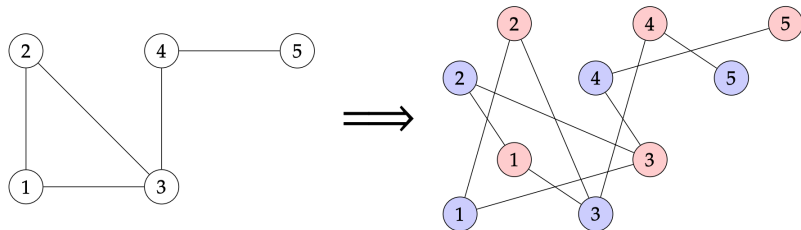
→ whether the length of a path  $P$  is even can be checked by comparing the colors of the start and end point

→ Example:  $P = 1, 2, 3, 1, 2$  has an even length because  $P_{\text{first}} = 1$  and  $P_{\text{last}} = 2$  have the same color.



## Trick: Graph Duplication

We modify our graph by duplicating every node and every edge:



→ whether the length of a path  $P$  is even can be checked by comparing the colors of the start and end point

→ Example:  $P = 1, 2, 3, 1, 2$  has an even length because  $P_{\text{first}} = 1$  and  $P_{\text{last}} = 2$  have the same color.

**Conclusion for the new graph**

A path between two nodes of the same color will always have even length

# Problem: Easteregg (SOI 2017)

## Easier problem

Find the shortest path  $P$  between  $a$  and  $b$  of *even* length.

# Problem: Easteregg (SOI 2017)

## Easier problem

Find the shortest path  $P$  between  $a$  and  $b$  of *even* length.

## Solution

- 1 Perform a graph duplication to get a new graph  $G'$ .

# Problem: Easteregg (SOI 2017)

## Easier problem

Find the shortest path  $P$  between  $a$  and  $b$  of *even* length.

## Solution

- 1 Perform a graph duplication to get a new graph  $G'$ .
- 2 Find the shortest path between  $a$  and  $b$  in the new graph  $G'$  (or between  $a$  and  $b$ ) using BFS.

# Easteregg: Lessons Learned

- Reducing a complex problem to an easier problem.
- Even if we are given a graph, we should not be afraid to modify it!
- Here we have seen one technique called "Graph Duplication".
- Other techniques:
  - duplicate graph more than two times
  - add dummy nodes
  - add dummy edges
  - ... be creative!

# Outline

- 1 Implicit Graphs
- 2 Graph Duplication
- 3 State Space Graphs**

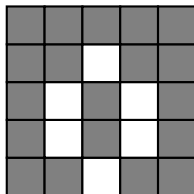
# Problem: Lights Out!

## Description

Given a 5x5 grid with some lights on or off. Clicking a cell toggles that cell and each of its immediate neighbors.

Goal: Turn off the lights with minimum number of clicks.

## Example



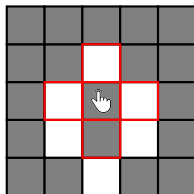
# Problem: Lights Out!

## Description

Given a 5x5 grid with some lights on or off. Clicking a cell toggles that cell and each of its immediate neighbors.

Goal: Turn off the lights with minimum number of clicks.

## Example





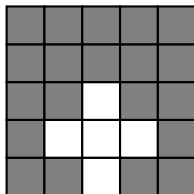
# Problem: Lights Out!

## Description

Given a 5x5 grid with some lights on or off. Clicking a cell toggles that cell and each of its immediate neighbors.

Goal: Turn off the lights with minimum number of clicks.

## Example



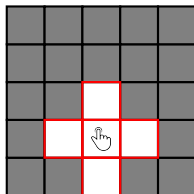
# Problem: Lights Out!

## Description

Given a 5x5 grid with some lights on or off. Clicking a cell toggles that cell and each of its immediate neighbors.

Goal: Turn off the lights with minimum number of clicks.

## Example



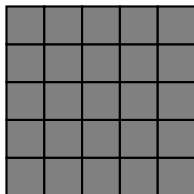
# Problem: Lights Out!

## Description

Given a 5x5 grid with some lights on or off. Clicking a cell toggles that cell and each of its immediate neighbors.

Goal: Turn off the lights with minimum number of clicks.

## Example



# State space

## Definition

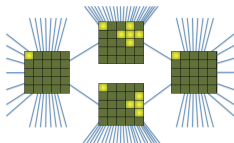
Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

# State space

## Definition

Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

## Example: Lights Out!



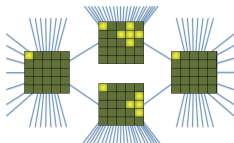
(Source: CMU)

# State space

## Definition

Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

## Example: Lights Out!



(Source: CMU)

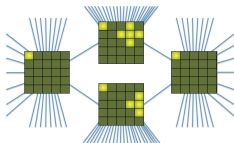
Q: How many states?

# State space

## Definition

Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

## Example: Lights Out!



(Source: CMU)

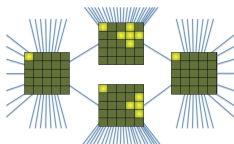
Q: How many states? A:  $2^{25} \approx 32\text{mio.}$

# State space

## Definition

Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

## Example: Lights Out!



(Source: CMU)

Q: How many states? A:  $2^{25} \approx 32\text{mio.}$

Q: How many edges?

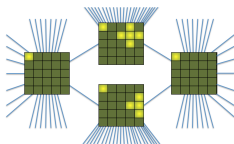


# State space

## Definition

Set of states that a problem can be in. The set of states forms a graph where two states are connected if there is an operation that can be performed to transform the first state into the second.

## Example: Lights Out!



(Source: CMU)

Q: How many states? A:  $2^{25} \approx 32$  mio.

Q: How many edges? A:  $25 \cdot 2^{24} \approx 400$  mio.

## Problem: Lights Out!

In order to avoid storing 32 mio. states and 400 mio. edges we will use an implicit graph!

## Problem: Lights Out!

In order to avoid storing 32 mio. states and 400 mio. edges we will use an implicit graph!

Implicitly retrieve the neighbor states of current

```
for(int i = 0; i < 5; i++) {
  for(int j = 0; j < 5; j++) {
    // current.click(i, j) simulates a click on (i, j)
    Board neighbor = current.click(i, j);
    // do something with neighbor
  }
}
```

## Problem: Lights Out!

In order to avoid storing 32 mio. states and 400 mio. edges we will use an implicit graph!

Implicitly retrieve the neighbor states of current

```
for(int i = 0; i < 5; i++) {
  for(int j = 0; j < 5; j++) {
    // current.click(i, j) simulates a click on (i, j)
    Board neighbor = current.click(i, j);
    // do something with neighbor
  }
}
```

## Solution

Perform a BFS on the state space graph (implicit graph)

# Lights Out: Lessons Learned

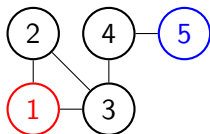
- Concepts of state space graph
- A very simple reduction between a problem into a space state graph problem

## Easter Egg (Revisited)

Instead of graph duplication we define a state space for the problem. We define the state space as (CurrentNode, CurrentLengthParity).

## Easter Egg (Revisited)

Instead of graph duplication we define a state space for the problem. We define the state space as (CurrentNode, CurrentLengthParity).



$$\text{StateSpace} = \{$$

$$(1, 0), (1, 1),$$

$$(2, 0), (2, 1),$$

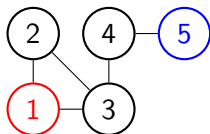
$$(3, 0), (3, 1),$$

$$(4, 0), (4, 1),$$

$$(5, 0), (5, 1)\}$$

## Easter Egg (Revisited)

Instead of graph duplication we define a state space for the problem. We define the state space as (CurrentNode, CurrentLengthParity).



$$\text{StateSpace} = \{$$

$$(1, 0), (1, 1),$$

$$(2, 0), (2, 1),$$

$$(3, 0), (3, 1),$$

$$(4, 0), (4, 1),$$

$$(5, 0), (5, 1)\}$$

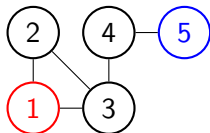
### Neighbor states of $(u, p)$

If  $(u, v)$  is an edge in the given graph, then  $(v, 1 - p)$  is a neighbor state of  $(u, p)$ .



## Easter Egg (Revisited)

Instead of graph duplication we define a state space for the problem. We define the state space as (CurrentNode, CurrentLengthParity).



$$\text{StateSpace} = \{$$

$$(1, 0), (1, 1),$$

$$(2, 0), (2, 1),$$

$$(3, 0), (3, 1),$$

$$(4, 0), (4, 1),$$

$$(5, 0), (5, 1)\}$$

### Neighbor states of $(u, p)$

If  $(u, v)$  is an edge in the given graph, then  $(v, 1 - p)$  is a neighbor state of  $(u, p)$ .

### Solution

Find shortest path from  $(a, 0)$  to  $(b, 0)$  in state space graph.

# Shortest Path of Even Length with even Sum

## Formal Statement

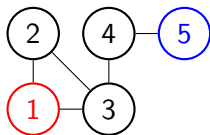
Find the shortest path  $P$  between  $a$  and  $b$ , such that the length of  $P$  is even and the sum of vertex numbers in  $P$  is even as well (i.e.  $\sum P_i$  is even).

# Shortest Path of Even Length with even Sum

## Formal Statement

Find the shortest path  $P$  between  $a$  and  $b$ , such that the length of  $P$  is even and the sum of vertex numbers in  $P$  is even as well (i.e.  $\sum P_i$  is even).

Use state: (CurrentNode, CurrentLengthParity, **CurrentSumParity**).



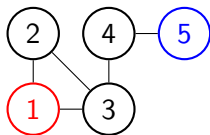
StateSpace = {  
 (1, 0, 0), (1, 0, 1), (1, 1, 0), (1, 0, 0),  
 (2, 0, 0), (2, 0, 1), (2, 1, 0), (2, 0, 0),  
 ...,  
 (5, 0, 0), (5, 0, 1), (5, 1, 0), (5, 0, 0)}

# Shortest Path of Even Length with even Sum

## Formal Statement

Find the shortest path  $P$  between  $a$  and  $b$ , such that the length of  $P$  is even and the sum of vertex numbers in  $P$  is even as well (i.e.  $\sum P_i$  is even).

Use state: (CurrentNode, CurrentLengthParity, **CurrentSumParity**).



StateSpace = {  
 (1, 0, 0), (1, 0, 1), (1, 1, 0), (1, 0, 0),  
 (2, 0, 0), (2, 0, 1), (2, 1, 0), (2, 0, 0),  
 ...,  
 (5, 0, 0), (5, 0, 1), (5, 1, 0), (5, 0, 0)}

## Neighbor states of $(u, p, q)$

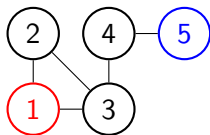
If  $(u, v)$  is an edge in the given graph, then  $(v, 1 - p, (v + q) \bmod 2)$  is a neighbor state of  $(u, p, q)$ .

# Shortest Path of Even Length with even Sum

## Formal Statement

Find the shortest path  $P$  between  $a$  and  $b$ , such that the length of  $P$  is even and the sum of vertex numbers in  $P$  is even as well (i.e.  $\sum P_i$  is even).

Use state: (CurrentNode, CurrentLengthParity, **CurrentSumParity**).



StateSpace = {  
 $(1, 0, 0), (1, 0, 1), (1, 1, 0), (1, 0, 0),$   
 $(2, 0, 0), (2, 0, 1), (2, 1, 0), (2, 0, 0),$   
 $\dots,$   
 $(5, 0, 0), (5, 0, 1), (5, 1, 0), (5, 0, 0)$ }

## Neighbor states of $(u, p, q)$

If  $(u, v)$  is an edge in the given graph, then  $(v, 1 - p, (v + q) \bmod 2)$  is a neighbor state of  $(u, p, q)$ .

Solution: Find shortest path from  $(a, 0, a \bmod 2)$  to  $(b, 0, 0)$ .

# Lessons Learned

- Using a state space graph to represent a graph problem.
- We can use multiple dimensions to represent our state.

# Coding Tasks

## Task 1: Farmland ( $\approx$ Maze)

Try to improve the code that implicitly retrieves the neighbors of  $(i, j)$ .

Hint: Use loops

## Task 2: Easteregg

Solve it two times:

- Duplicate the graph explicitly and perform a BFS.
- Use the state (node, parity) and use implicit graphs.

## Task 3: Lights Out! (Optional)

Hint: Use a 32 bit integer to represent the state.